

Counting distributions in Bayesian networks and Markov models with polynomials

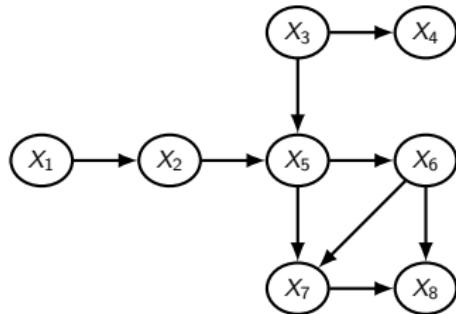
Alexandra Lefebvre and Grégory Nuel

Laboratoire de Probabilités Statistique et Modélisation, CNRS, Sorbonne Université, Paris

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Bayesian networks - DAG - and belief propagation



$$X_{\mathcal{U}} = \{X_u\}_{u \in \mathcal{U}}$$
$$\forall u \in \mathcal{U}, X_u \in \mathcal{X}_u$$

pa_u : set of labels for parents of X_u
 $\sum_{X_{\mathcal{U}}} \prod_{u \in \mathcal{U}} \mathbb{P}(X_u | X_{\text{pa}_u}) = 1$

Let $\overline{X_u} = \{X_u, X_{\text{pa}_u}\}$,

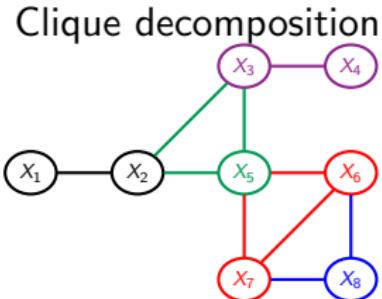
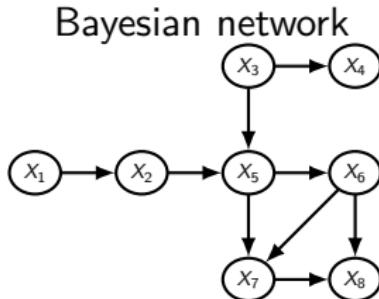
$$\text{ev} = \{X_u \in \mathcal{X}_u^* \subset \mathcal{X}_u\} \quad \rightarrow \quad \varphi_{X_u}(\overline{X_u}) = \mathbb{1}_{X_u \in \mathcal{X}_u^*} \mathbb{P}(X_u | X_{\text{pa}_u})$$

$$\mathbb{P}(\text{ev}) = \sum_{X_{\mathcal{U}}} \prod_{u \in \mathcal{U}} \varphi_{X_u}(\overline{X_u})$$

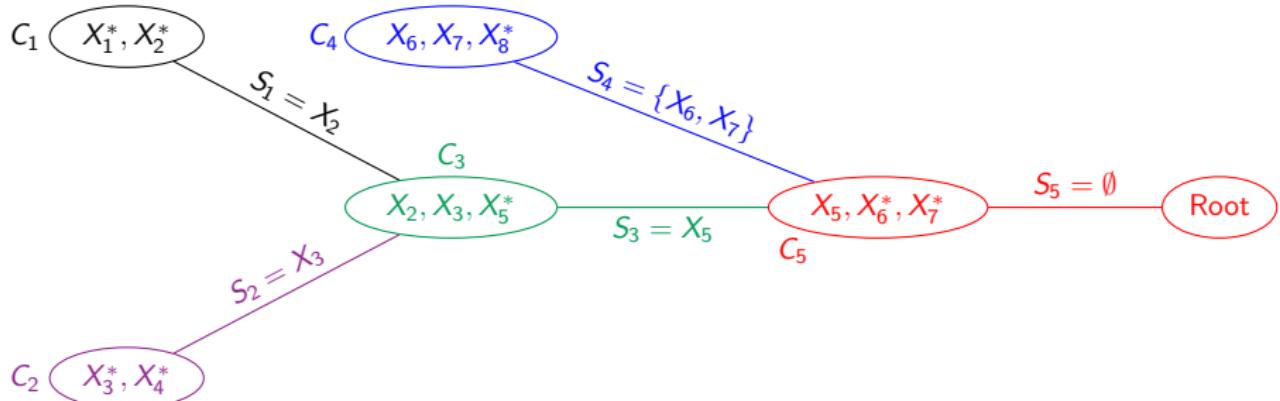
$$\forall v \in \mathcal{U}, \quad \mathbb{P}(X_v, \text{ev}) = \sum_{X_{\mathcal{U}} \setminus X_v} \prod_{u \in \mathcal{U}} \varphi_{X_u}(\overline{X_u})$$

Complexity = $\mathcal{O}(\prod_u |\mathcal{X}_u|)$, ie $\mathcal{O}(2^n)$ for only binary variables.

Message passing or sum-product algorithm



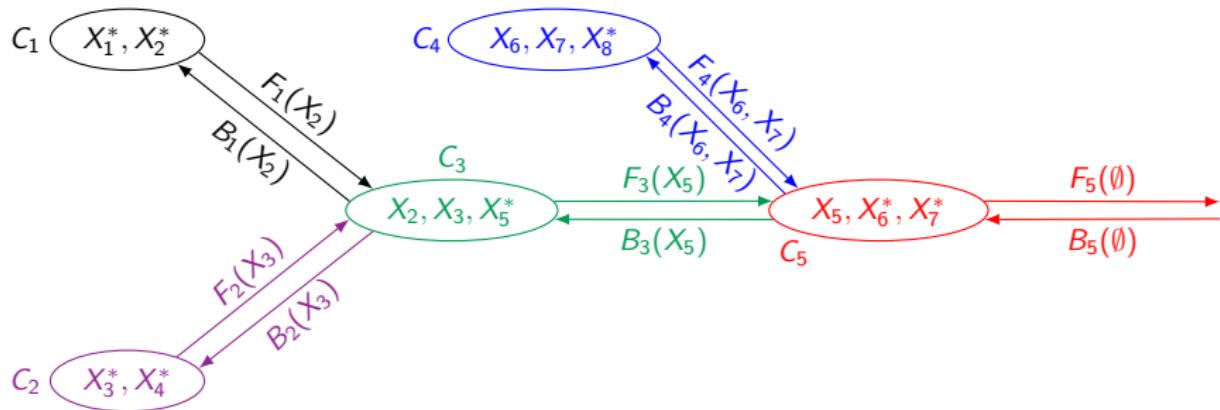
$$\begin{aligned} \mathbb{P}(\text{ev}) = & \sum_{X_5} \sum_{X_6} \sum_{X_7} \left\{ \sum_{X_8} \varphi_{X_8}(\bar{X}_8) \right. \\ & \left. \sum_{X_2} \sum_{X_3} \left(\underbrace{\sum_{X_1} \varphi_{X_1}(\bar{X}_1) \varphi_{X_2}(\bar{X}_2)}_{F_1(X_2)} \underbrace{\sum_{X_4} \varphi_{X_3}(\bar{X}_3) \varphi_{X_4}(\bar{X}_4) \varphi_{X_5}(\bar{X}_5)}_{F_2(X_3)} \right. \right. \\ & \left. \left. \varphi_{X_6}(\bar{X}_6) \varphi_{X_7}(\bar{X}_7) \right\} \right. \end{aligned}$$



- Junction-tree properties :

- Tree
- Running intersection
- Covering

Message passing



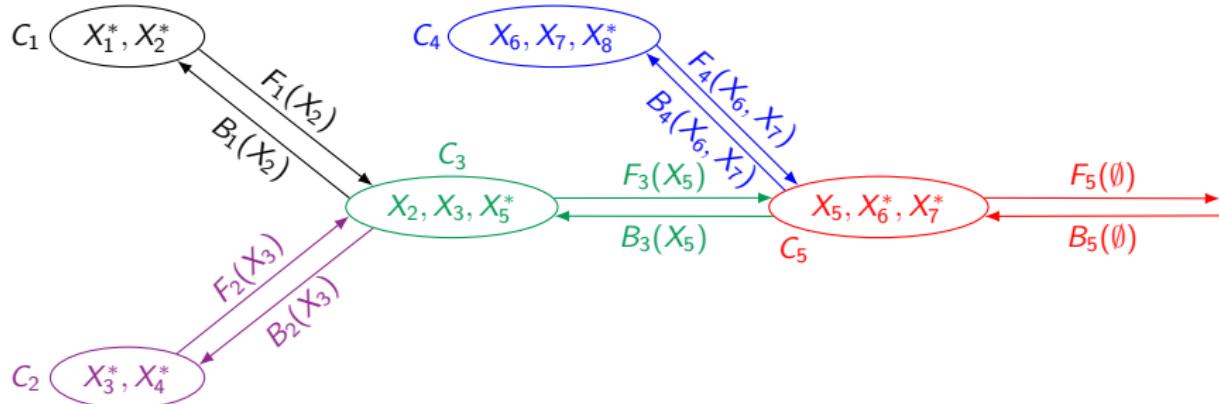
$$\forall i \in \{1, \dots, K\}; \quad \phi_i(C_i) = \prod_{X_u \in C_i^*} \varphi_{X_u}(\overline{X_u})$$

$$F_i(S_i) = \sum_{C_i \setminus S_i} \left(\prod_{j \in \text{from}_i} F_j(S_j) \right) \times \phi_i(C_i)$$

$$B_K(\emptyset) = 1 \text{ and for } k \in K, \dots, 1, \forall i \in \text{from}_k,$$

$$B_i(S_i) = \sum_{C_k \setminus S_i} \phi_k(S_k) \times B_k(S_k) \times \prod_{j \in \text{from}_k, j \neq i} F_j(S_j)$$

Message passing



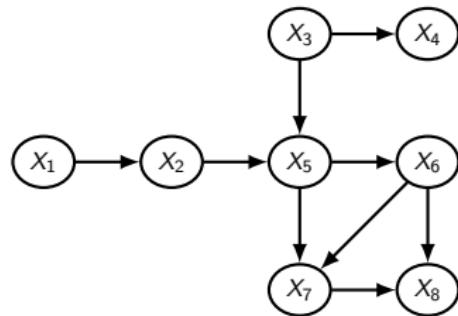
$\forall i \in \{1, \dots, K\}$,

$$\mathbb{P}(S_i, \text{ev}) = F_i(S_i) \times B_i(S_i)$$

$$\mathbb{P}(C_i, \text{ev}) = \phi_i(C_i) \times B_i(S_i) \times \prod_{j \in \text{from}_i, j \neq i} F_j(S_j)$$

Complexity = $\mathcal{O}(n \times \text{argmax}_u |\mathcal{X}_u|^{\text{TW}})$, TW = size of the largest clique

Idea : Change definition of ev and work on the arithmetic of potentials to compute other quantities of interest.



$$X_{\mathcal{U}} = \{X_u\}_{u \in \mathcal{U}}$$
$$\forall u \in \mathcal{U}, X_u \in \mathcal{X}_u$$

$$\text{pa}_u : \text{set of labels for parents of } X_u$$
$$\sum_{X_{\mathcal{U}}} \prod_{u \in \mathcal{U}} \mathbb{P}(X_u | X_{\text{pa}_u}) = 1$$

Let $\overline{X_u} = \{X_u, X_{\text{pa}_u}\}$

$$\text{ev} = \{X_u \in \mathcal{X}_u^* \subset \mathcal{X}_u\} \quad \rightarrow \quad \varphi_{X_u}(\overline{X_u}) = \mathbb{1}_{X_u \in \mathcal{X}_u^*} \mathbb{P}(X_u | X_{\text{pa}_u})$$

$$\mathbb{P}(\text{ev}) = \sum_{X_{\mathcal{U}}} \prod_{u \in \mathcal{U}} \varphi_{X_u}(\overline{X_u})$$

$$\forall v \in \mathcal{U}, \quad \mathbb{P}(X_v, \text{ev}) = \sum_{X_{\mathcal{U}} \setminus X_v} \prod_{u \in \mathcal{U}} \varphi_{X_u}(\overline{X_u})$$

Derivatives of the likelihood with derivative generating functions in parametric bayesian networks.

$$\varphi_u(\overline{X_u}) = \mathbb{1}_{X_u \in \mathcal{X}_u^*} \mathbb{P}(X_u | X_{\text{pa}_u}, \theta) \quad \Rightarrow \quad \sum_{X_u} \prod_{u \in \mathcal{U}} \varphi_u(\overline{X_u}) = \mathbb{P}(\text{ev} | \theta) = L(\theta)$$

Message passing algo.: $\forall i \in 1, \dots, K, \quad \sum_{S_i} F_i(S_i) B_i(S_i) = L(\theta)$

with K : Number of cliques of the built JT.

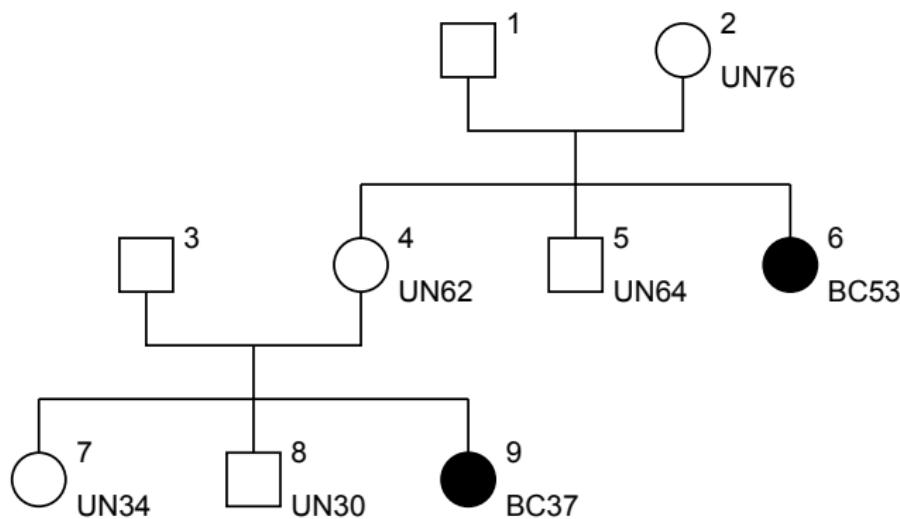
Polynomial potentials: $\Phi_u(\overline{X_u}) = \sum_{k=0}^d \varphi_u^{(d)}(\overline{X_u}) z^d$

Leibniz's product: $\sum_{k=0}^d a_k z^k \star \sum_{k=0}^d b_k z^k = \sum_{k=0}^d \sum_{i=0}^k \binom{i}{k} a_i b_{k-i} z^k$

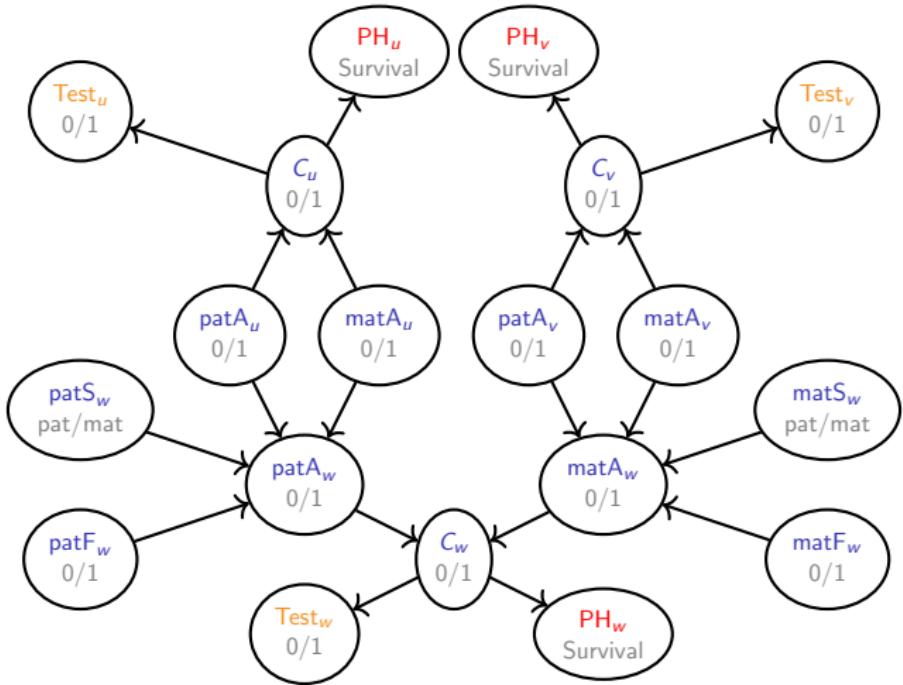
Message passing algo over polynomial potentials:

$\forall i \in 1, \dots, K, \quad \sum_{S_i} \mathbf{F}_i(S_i) \mathbf{B}_i(S_i) = \sum_{k=0}^d L^{(k)}(\theta) z^k$

Counting distribution with probability and moment generating functions (Application with genetic diseases).



- Distribution of the # of carriers (N^C)
- Distribution of the # of filiation errors (N^F) (+joint distributions)
- Moments



$$\text{ev} = \{\text{PH}_u\}_{u=1,\dots,n_{\text{ind}}} \cap \{\text{Test}_u\}_{u \in \{1,\dots,n_{\text{ind}}\}}, \quad n_{\text{ind}} = \# \text{ of ind.}$$

Probability generating function of N : $G_N(z) = \mathbb{E}(z^N)$
 with $N = \sum_u C_u$ (Number of carriers).

Evidence: $\text{ev} = \{\text{PH}_u\}_{u=1,\dots,n_{\text{ind}}} \cap \{\text{Test}_u\}_{u \in \{1,\dots,n_{\text{ind}}\}}$

Polynomial Potentials: $\varphi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u)$
 $\Rightarrow \Phi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u) z^{C_u}$

For each other variable, $\varphi_{\cdot}(\cdot) = \Phi_{\cdot}(\cdot)$

Computation with the Sum-product / message passing algorithm:
 With variables renamed X_u , $u \in \mathcal{U} = \{1, \dots, n\}$ ($n = \#$ of variables):
 $\forall i \in 1, \dots K$ (K : # of cliques of the built JT),

$$\begin{aligned} \sum_s \underbrace{\mathbf{F}_i(s) \times \mathbf{B}_i(s)}_{\text{Polynomial messages}} &= \sum_{X_u} \prod_u \Phi_{X_u}(\overline{X_u}) \\ &= \sum_{k \geq 0} \mathbb{P}(N = k, \text{ev}) z^k \end{aligned}$$

multivariate polynomials for joint distributions

$$N^C = \sum_u C_u \text{ and } N^F = \sum_u \text{patF}_u$$

Joint distribution of # of carriers and # of paternal filiation errors

For all individuals: $\varphi_{C_u}(\overline{C_u}) \Rightarrow \Phi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u) z^{C_u}$

For offsprings: $\varphi_{\text{patF}_u}(\overline{\text{patF}_u}) \Rightarrow \Phi_{\text{patF}_u}(\overline{\text{patF}_u}) = \mathbb{P}(\text{patF}_u) y^{\text{patF}_u}$

Computation with the sum-product / message-passing algorithm:

With variables renamed X_u , $u \in \mathcal{U} = \{1, \dots, n\}$ ($n = \#$ of variables):

$\forall i \in 1, \dots K$ (K : # of cliques of the built JT),

$$\begin{aligned} \sum_s \underbrace{\mathbf{F}_i(s) \times \mathbf{B}_i(s)}_{\text{Polynomial messages}} &= \sum_{X_U} \prod_u \Phi_{X_u}(\overline{X_u}) \\ &= \sum_{k \geq 0} \sum_{\ell \geq 0} \mathbb{P}(N^C = k, N^F = \ell, \text{ev}) z^k y^\ell \end{aligned}$$

Moment generating function of N : $M_N(t) = \mathbb{E}(e^{tN})$

Polynomial potentials: $\varphi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u)$

$$\Rightarrow \Phi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u) e^{t C_u}$$

$$\Rightarrow \Phi_{C_u}(\overline{C_u}) = \mathbb{P}(C_u | \text{patA}_u, \text{matA}_u) \left[\sum_{\ell=0}^d \frac{t^\ell}{\ell!} \right]^{C_u} \text{ with } d, \text{ a chosen order.}$$

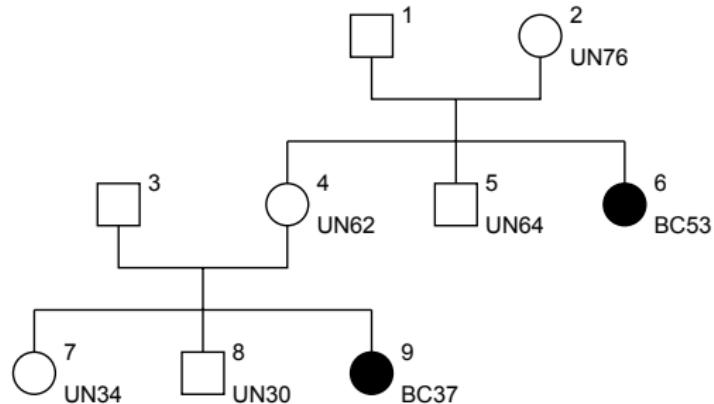
For each other variable, $\varphi_{\cdot}(\cdot) = \Phi_{\cdot}(\cdot)$

Computation with the sum-product / message-passing algorithm:

$\forall i \in 1, \dots, K$,

$$\begin{aligned} \sum_s \underbrace{\mathbf{F}_i(s) \times \mathbf{B}_i(s)}_{\text{Polynomial messages}} &= \sum_{m \geq 0} \mathbb{P}(N = m, \text{ev}) e^{mt} \\ &= \sum_{m \geq 0} \sum_{\ell=0}^d \mathbb{P}(N = m, \text{ev}) \frac{m^\ell t^\ell}{\ell!} = \sum_{\ell=0}^d \underbrace{\left\{ \sum_{m \geq 0} \mathbb{P}(N = m, \text{ev}) m^\ell \right\}}_{\mathbb{E}[N^\ell, \text{ev}]} \frac{t^\ell}{\ell!} \end{aligned}$$

Results



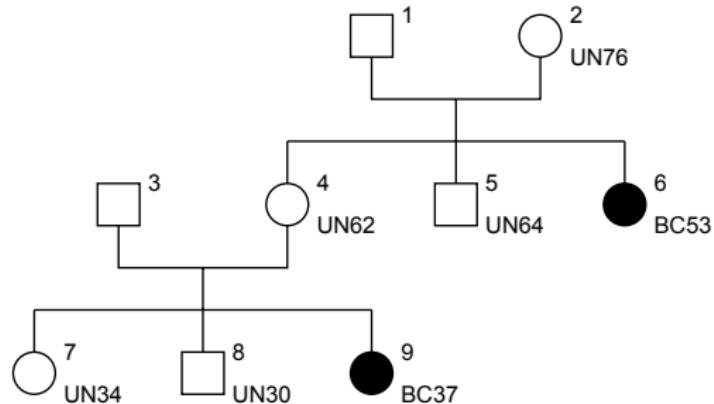
$$\pi(\cdot) = \mathbb{P}(\cdot | \text{FH})$$

$2/3$	$2/4$	$2/6$	$2/9$	$\pi(C_J)$
0	0	0	0	67.5
0	1	1	1	18.0
1	0	0	1	10.0
0	0	1	0	1.9
0	1	0	1	1.4
0	1	1	0	0.4

Individual i	$\pi(N)$	2/1	2/2	2/3	2/4	2/5	2/6	2/7	2/8	2/9
$\pi(C_i = 1)$	—	17.3	5.1	10.7	20.0	11.1	20.8	14.7	15.2	30.0
$\pi(C_i = 1 N = 0)$	67.3	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0
$\pi(C_i = 1 N = 5)$	7.6	77.6	22.4	1.6	98.5	36.9	93.1	34.7	37.0	98.1
$\pi(C_i = 1 N = 6)$	7.0	77.6	22.6	2.1	98.5	68.1	97.5	66.0	68.3	99.3
$\pi(C_i = 1 N = 3)$	6.2	14.6	4.2	81.1	4.1	14.8	15.6	39.7	42.5	83.5
$\pi(C_i = 1 N = 4)$	5.5	44.0	12.6	44.2	55.7	4.4	46.6	47.5	47.8	97.1
$\pi(C_i = 1 N = 2)$	3.8	20.9	5.9	73.1	0.3	1.9	24.8	1.4	1.5	70.4

3 founders, 6 offsprings, 60 variables, 41 cliques, complexity 748

Results



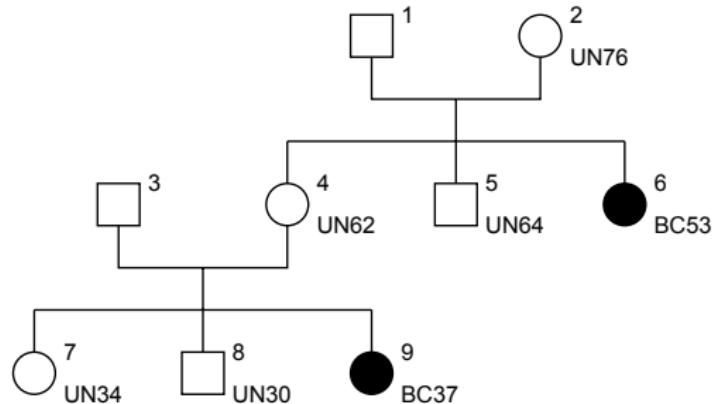
$$\pi(\cdot) = \mathbb{P}(\cdot | \text{FH}, T_4 = 1)$$

2/3	2/4	2/6	2/9	$\pi(C_J)$
0	1	1	1	81.8
0	0	0	0	7.7
0	1	0	1	6.4
0	1	1	0	1.7
1	0	0	1	1.1
1	1	1	1	0.8

Individual i	$\pi(N)$	2/1	2/2	2/3	2/4	2/5	2/6	2/7	2/8	2/9
$\pi(C_i = 1)$	—	70.8	20.7	2.0	90.9	45.4	84.6	44.6	46.2	90.2
$\pi(C_i = 1 N = 5)$	34.1	77.6	22.4	0.2	100	36.9	93.1	34.7	37.0	98.1
$\pi(C_i = 1 N = 6)$	31.3	77.6	22.6	0.7	100	68.0	97.5	66.0	68.3	99.3
$\pi(C_i = 1 N = 4)$	14.1	76.2	21.8	2.0	98.1	7.6	80.5	9.2	9.7	95.0
$\pi(C_i = 1 N = 7)$	10.2	77.7	23.6	3.8	100	97.4	99.8	98.9	99.0	100
$\pi(C_i = 1 N = 0)$	7.6	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0
$\pi(C_i = 1 N = 3)$	1.8	51.5	14.6	31.2	63.1	6.6	18.0	17.4	18.7	78.9

3 founders, 6 offsprings, 60 variables, 41 cliques, complexity 748

Results



$$\pi(\cdot) = \mathbb{P}(\cdot | \text{FH}, T_4 = 0)$$

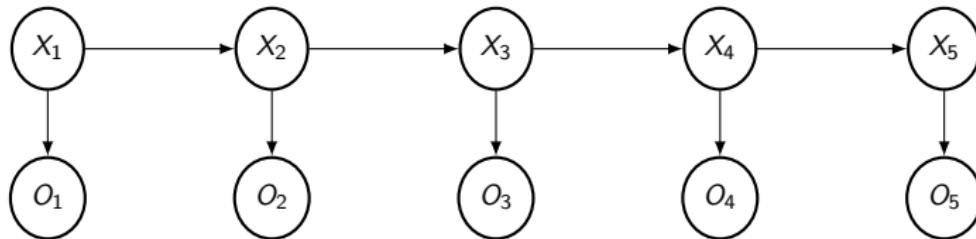
2/3	2/4	2/6	2/9	$\pi(C_J)$
0	0	0	0	80.3
1	0	0	1	11.9
0	1	1	1	4.4
0	0	1	0	2.2
0	1	0	1	0.3
1	0	1	1	0.3

Individual i	$\pi(N)$	2/1	2/2	2/3	2/4	2/5	2/6	2/7	2/8	2/9
$\pi(C_i = 1)$	—	5.9	1.7	12.5	4.9	3.8	7.1	8.3	8.6	17.1
$\pi(C_i = 1 N = 0)$	80.1	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0
$\pi(C_i = 1 N = 3)$	7.1	12.6	3.6	83.9	0.9	15.2	15.5	40.9	43.8	83.7
$\pi(C_i = 1 N = 2)$	4.5	20.9	5.9	73.2	0.1	1.9	24.9	1.4	1.5	70.4
$\pi(C_i = 1 N = 4)$	3.6	17.2	4.9	79.5	20.4	1.8	18.2	79.5	79.6	98.8
$\pi(C_i = 1 N = 5)$	2.0	77.6	22.4	7.2	93.0	36.9	93.1	34.8	37.1	98.1
$\pi(C_i = 1 N = 6)$	1.8	77.6	22.6	7.5	93.1	68.1	97.5	66.0	68.3	99.3

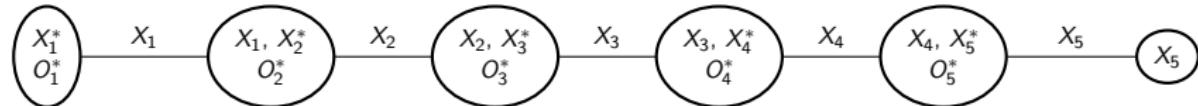
3 founders, 6 offsprings, 60 variables, 41 cliques, complexity 748

Number of errors in a Markov Chain, $N = \sum_i \mathbb{1}_{X_i \neq O_i}$

Hidden Markov chain



Junction tree



Potentials

$$\varphi_{X_1}(X_1) = \mu_1; \varphi_{X_i}(\bar{X}_i) = \pi(X_{i-1}, X_i) \text{ with } \pi = \begin{pmatrix} 0.9 & 0.1 \\ 0.2 & 0.8 \end{pmatrix}$$

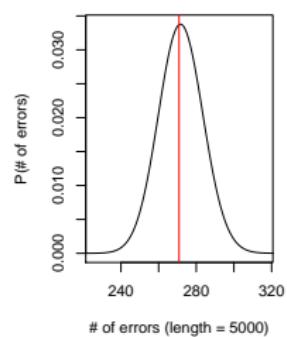
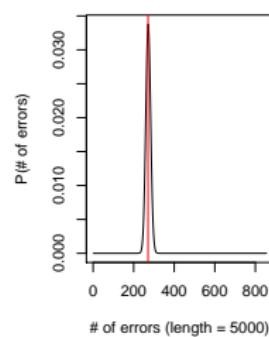
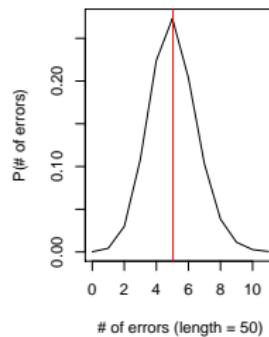
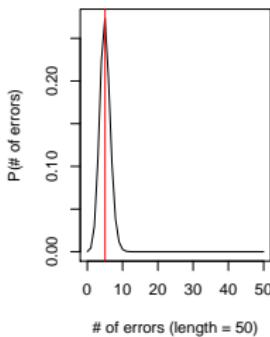
$$\text{PGF: } \varphi_{O_i}(O_i, X_i = j) = \underbrace{\mathbb{1}_{O_i = \omega_i}}_{\text{ev}} (1 - \eta)^{\mathbb{1}_{O_i=j}} (\eta \times z)^{\mathbb{1}_{O_i \neq j}} \text{ with } \eta = 0.05$$

$$\text{MGF: } \varphi_{O_i}(O_i, X_i = j) = \underbrace{\mathbb{1}_{O_i = \omega_i}}_{\text{ev}} (1 - \eta)^{\mathbb{1}_{O_i=j}} (\eta(1 + t + t^2/2 + \dots))^{\mathbb{1}_{O_i \neq j}}$$

$$\text{PGF: } \forall i \in 1, \dots K, \quad \sum_{X_i} \mathbf{F}_i(X_i) \mathbf{B}_i(X_i) = \sum_{k \geq 0} \mathbb{P}(N = k, \text{ev}) z^k$$

$$\text{MGF: } \forall i \in 1, \dots K, \quad \sum_{X_i} \mathbf{F}_i(X_i) \mathbf{B}_i(X_i) = \sum_{\ell \geq 0} \mathbb{E}[N^\ell, \text{ev}] \frac{t^\ell}{\ell!}$$

Distributions:



Moments:

$$m_0 = 1.00, \quad m_1 = 5.04, \quad m_2 = 27.63$$

$$m_0 = 1.00, \quad m_1 = 270.85, \quad m_2 = 73501.20$$

Number of atypical segments in a sequence

Joint work with S. Mercier

Observed variables: $X = X_1, \dots, X_n \in \mathcal{X}^n$ (e.g. $\mathcal{X} = \{\text{A, C, G, T}\}$)

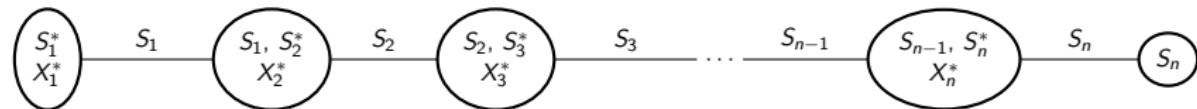
Latent variables: $S = S_1, \dots, S_n \in \mathcal{S}^n$ ($\mathcal{S} = \{1, 2, 3\}$)

Emission: $\mathbb{P}(X_i|S_i) = \begin{cases} q_0(X_i) & \text{if } S_i \neq 2 \\ q_1(X_i) & \text{if } S_i = 2 \end{cases}$ Constr. transition $\pi = \begin{pmatrix} 110 \\ 011 \\ 001 \end{pmatrix}$

DAG



Junction tree



- **Evidence:** $\forall x \in \mathcal{X}$, $e(x) = \frac{q_1(x)}{q_0(x)}$ or $e(x) = \exp\left(\frac{f(x)}{T}\right)$

- **Potentials:**

$$\phi_1(S_1 = k) = \pi(1, k)e(x_1)^{\mathbb{1}_{k=2}}$$

$$\text{For } i = 2, \dots, n \quad \phi_i(S_{i-1} = j, S_i = k) = \pi(j, k)e(x_i)^{\mathbb{1}_{k=2}}$$

- **Forward - Backward:**

$$F_1(S_1) = \phi_1(S_1) \quad \text{and for } i = 2, \dots, n \quad F_i(k) = \sum_j F_{i-1}(j)\phi_i(j, k)$$

$$B_n(\emptyset) = 1 \text{ (conv.)} \quad \text{and for } i = n, \dots, 2, \quad B_{i-1}(j) = \sum_k \phi(j, k)B_i(k)$$

- **Start and stop of a segment**

$$\text{P(Start at pos. } i\text{): } \mathbb{P}(S_i = 2, S_{i-1} = 1, X) \propto F_{i-1}(1)\phi_i(1, 2)B_i(2)$$

$$\text{P(Stop at pos. } i\text{): } \mathbb{P}(S_i = 2, S_{i+1} = 3, X) \propto F_i(2)\phi_i(2, 3)B_{i+1}(3)$$

- **Probability of the sequence** $\forall i \in \{1, \dots, n\} \quad \mathbb{P}(X) \propto \sum_k F_i(k)B_i(k)$

N segments

First approach

$\mathcal{S} = \{1, 2, 3, 4, 5, \dots, 2N, 2N+1\}$ and π of dimension $(2N+1) \times (2N+1)$

With polynomials, pgf & mgf: $\mathcal{S} = \{0, 1\}$ and $\pi = \begin{pmatrix} 1 & z \\ 1 & 1 \end{pmatrix}$

$$\phi_1(S_1) = e(x_1)^{\mathbb{1}_{S_1}} \pi(S_0, S_1) \quad \text{with } S_0 = 0 \text{ (conv.)}$$

$$\forall i = 2, \dots, n \quad \phi_i(j, k) = e(x_i)^{\mathbb{1}_{S_i}} \pi(j, k)$$

Forward-Backward algo →

$$\forall i \in \{1, \dots, n\} \quad F_i(S_i) B_i(S_i) \propto \sum_k \mathbb{Q}(S_i, N=k, X) z^k$$

$$\text{and } \mathbb{P}(S_i, N=k, X) \propto \mathbb{Q}(S_i, N=k, X) / \binom{n+1}{2k}$$

Start / Stop positions with $N \geq 0$, Distribution of N

$$\mathbb{Q}(N = k, X) \propto \left[\sum_j F_i(j) B_i(j) \right]_{z^k}$$

and

$$\mathbb{P}(X) \propto \sum_k \mathbb{Q}(N = k, X)$$

$\mathbb{P}(\text{any segment starts at position } i)$

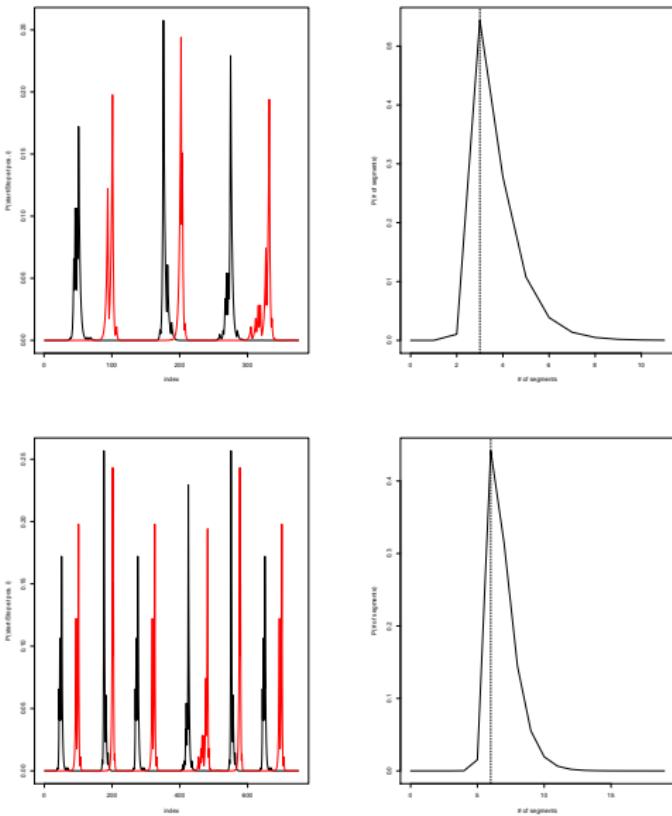
$$\mathbb{P}(S_{i-1} = 0, S_i = 1 | N = k, X) = [F_{i-1}(0) B_i(1) e(X_i) \times z]_{z^k} / \mathbb{Q}(N = k, X)$$

$\mathbb{P}(\text{any segment stops at position } i)$

$$\mathbb{P}(S_i = 1, S_{i+1} = 0 | N = k, X) = [F_i(1) B_{i+1}(0)]_{z^k} / \mathbb{Q}(N = k, X)$$

Distribution of N

$$\mathbb{P}(N = k | X) = \frac{\mathbb{Q}(N = k, X)}{\binom{n+1}{2k} \times \mathbb{P}(X)}$$



top : simu. with 3 atypical seg. bottom: simu. with 6 atypical seg.